

DOL appendix E: conformance of UML class diagrams

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Conformance of UML class diagrams to DOL means to provide an institution, i.e.

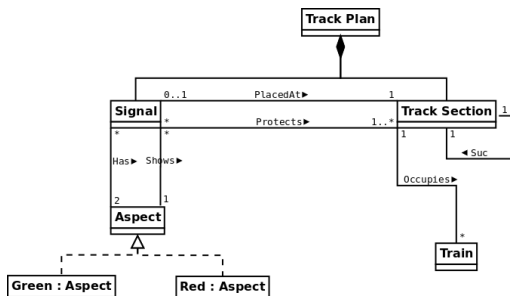
- **Signatures** (compact representation of the non-logical symbols occurring in the abstract syntax tree)
- **Models** (here, we do this via a translation to Common Logic)
- **Sentences** (compact representation of the constraints occurring in the abstract syntax tree)
- **Satisfaction** (of a sentence in a model)
- **Signature morphisms** (transcends proper UML) and associated model reducts and sentence translations

Signatures $\Sigma \in |\text{Sig}|$ comprise

- **Classifier hierarchy** (C, \leq)
e.g. `Flight < Travel`
- **Instance specifications** $k : c$
e.g. `LH123:Flight`
(OK, this is for object diagrams...)
- **Property declarations** $c.p(x_1 : c_1, \dots, x_n : c_n) : c'$
e.g. `Flight.number:Int`
- **Composition declarations** $c \blacklozenge r : c'$
e.g. `TrackPlan \blacklozenge signal : Signal`
- **Association declarations** $a(r_1 : c_1, \dots, r_n : c_n)$ e.g.
`PlacedAt(signal : Signal, trackSection : TrackSection)`

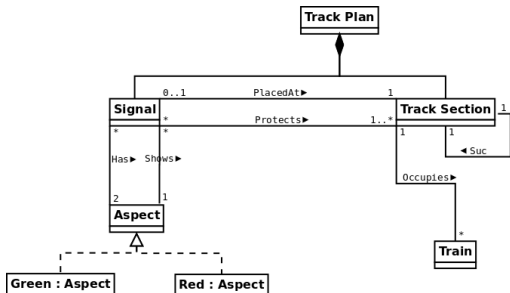
Signature morphisms map these in a compatible way.

Signature Example: Circle DSL (1)



- Classifiers: TrackPlan, Signal, TrackSection, Aspect, Train

Signature Example: Circle DSL (2)

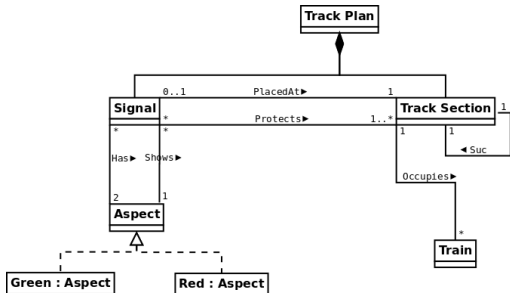


- Compositions:

TrackPlan◆signal : Signal,

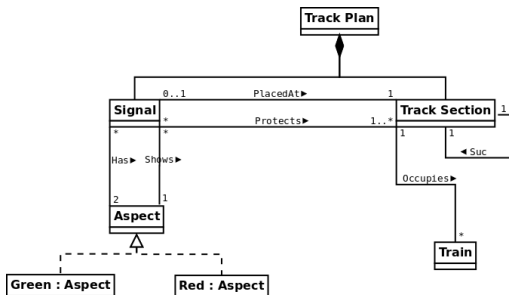
TrackPlan◆trackSection : TrackSection

Signature Example: Circle DSL (3)



- Associations, e.g.:
`PlacedAt({signal : Signal, trackSection : TrackSection}),`
`Protects({signal : Signal, trackSection : TrackSection}),`
`Suc({trackSection1 : TrackSection, trackSection2 :`
`TrackSection})`

Signature Example: Circle DSL (4)



- Instance specifications:
Green : Aspect, Red : Aspect

Models via translation to Common Logic

Models are inherited from Common Logic, via a translation:

For a class net $\Sigma = ((C, \leq_C), K, P, M, A)$, we define a Common Logic theory $CL(\Sigma)$ consisting of:

- for each **class** $c \in C$, a predicate $CL(c)$, such that
 - $CL(\text{Boolean}) = \text{buml:Boolean}$,
 - $CL(\text{String}) = \text{buml:String}$,
 - $CL(\text{Integer}) = \text{buml:Integer}$,
 - $CL(\text{UnlimitedNatural}) = \text{form:NaturalNumber}$,
 - $CL(\text{Real}) = \text{buml:Real}$,
 - $CL(\text{List}[c]) = \text{form:Sequence}$ — typed sequences???
 - $CL(\text{Set}[c]) = ???$
 - $CL(\text{Bag}[c]) = ???$
 - $CL(\text{Pair}[c_1, c_2]) = ???$
 - $CL(\text{Enumeration}[v_1, v_2, \dots, v_n]) = ???$
- for each **subclass relation** $c_1 \leq_C c_2$, an axiom $(\text{forall } (x) (\text{if } (C1 x) (C2 x)))$,
where $C1 = CL(c_1)$, $C2 = CL(c_2)$,

Models via translation to Common Logic (cont'd)

- CL maps each **instance specification** declaration $k : c \in K$ to constant $CL(k)$ and an axiom $(c \text{ k})$, where by abuse of notation, we identify c with $CL(c)$, and k with $(CL(k))$ (this abuse of notation will also be used in the sequel);
- for two instance specifications $k_1 : c$ and $k_2 : c$ with $k_1 \neq k_2$, an axiom $(\text{not } (= k_1 k_2))$ (the unique name assumption);
- CL maps each **property declaration** $c.p(x_1 : c_1, \dots, x_n : c_n) : c' \in P$ to a predicate $CL(c.p)$ and axioms
 - $(\text{forall } (x_1 \ x_2 \ \dots \ x_n \ x) (\text{if } (c.p \ x_1 \ x_2 \ \dots \ x_n \ x) (c_i \ x_i)))$ for each $i = 1 \dots n$,¹
 - $(\text{forall } (x_1 \ x_2 \ \dots \ x_n \ x) (\text{if } (c.p \ x_1 \ x_2 \ \dots \ x_n \ x) (c' \ x)))$
 - $(\text{forall } (x_1 \ x_2 \ \dots \ x_n \ x \ y) (\text{if } (\text{and } (c.p \ x_1 \ x_2 \ \dots \ x_n \ x) (c.p \ x_1 \ x_2 \ \dots \ x_n \ x)) (= x y)))$

¹Note that the \dots here is meta notation, not a sequence marker!

Models via translation to Common Logic (cont'd)

- CL maps each **composition declaration** $c \blacklozenge r : c' \in M$ to a predicate $CL(r)$ and an axiom
(forall (x y) (if (r x y) (and (c x) (c' y))))
- for any pair of composition declarations $c_1 \blacklozenge r_1 : c'_1$ and $c_2 \blacklozenge r_2 : c'_2$, an axiom stating “each instance has at most one owner”:
(forall (x₁ x₂ y) (if (and (r₁ x₁ y) (r₂ x₂ y)) (= x₁ x₂)))
- CL maps each **association declaration** $a(r_1 : c_1, \dots, r_n : c_n) \in A$ to a predicate $CL(a)$ and an axiom
(forall (x₁ x₂ ... x_n) (if (a x₁ x₂ ... x_n) (and (c₁ x₁) ... (c_n x_n))))

Sentences and Satisfaction

Sentences $\varphi \in |\text{Sen}(\Sigma)|$ capture *multiplicities*

- Comparing cardinality's $e \leq l$, $e \geq l$ (l a natural number)

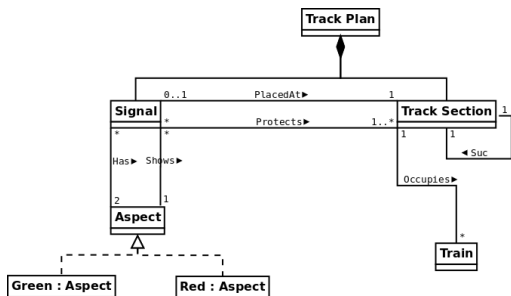
Cardinality expressions e

- ... of compositions
 - How many owned instances? $\#(c \blacklozenge r : c')$
 - For each $o \in c^{\circ}$, cardinality of $(c \blacklozenge r : c')^{\circ}(o)$
 - How many owners? $\#(c \blacklozenge r : c')[\text{owner}]$
- ... of associations
 - How many tuples, when fixing a subset of roles?
 $\#z(\{r_1 : c_1, \dots, r_n : c_n\})[r_{i_1}, \dots, r_{i_m}]$

Satisfaction relation $\mathcal{M} \models_{\Sigma} e \leq l$, $\mathcal{M} \models_{\Sigma} l \leq e$

- Compare each cardinality resulting from evaluating e to l

Sentences Example: Circle DSL



- Cardinality of associations, e.g. :
`#PlacedAt({signal : Signal, trackSection : TrackSection})[signal] ≤ 1`
`#PlacedAt({signal : Signal, trackSection : TrackSection})[signal] ≥ 1`